

(Maybe) Owning Things, Declaratively

Modelling Rust's Ownership in Datalog

Who?

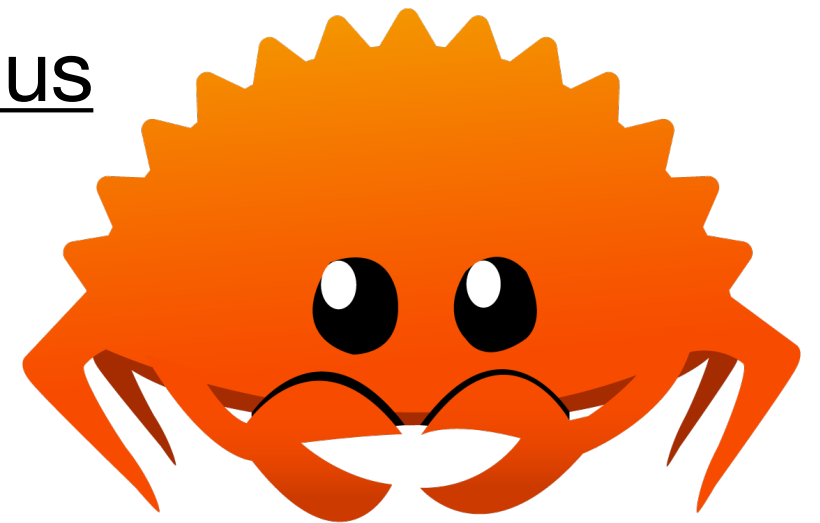
- I am Albin, and this is my Master's project
- Working with Niko Matsakis (et. al.) of Mozilla Research
- Tobias is subject reviewing



What?



- **Polonius**: a declarative reformulation of the [Rust borrow check](#)
- Separate package (“crate”) consumed by Rust
- Developed on GitHub: [rust-lang/polonius](https://github.com/rust-lang/polonius)
- Current status: *very* work-in-progress!



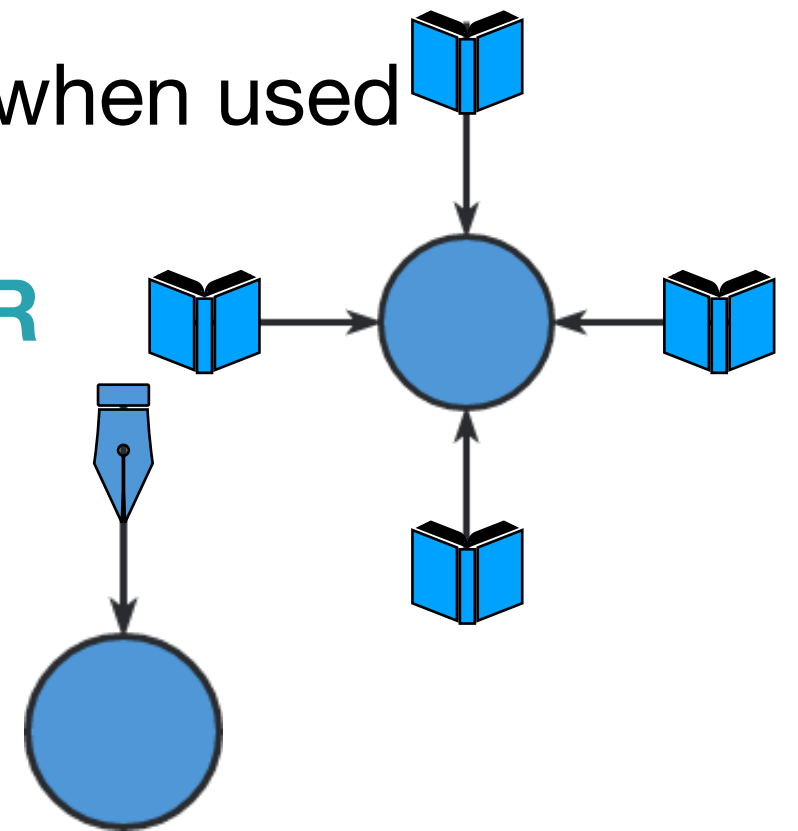
What??

- The borrow check: static analysis of references
- Used to ensure safe shared memory use
- May-point-to (“provenance”) part of a reference’s type
 - How do we interpret such a type? Stay tuned!

What???

House rules

- A value must be (initialised | not moved) when used
- N (live) read-references to same value **OR**
- ≤ 1 (live) write-reference to same value
- A reference must never outlive its value



What????

```
let mut x = 22;  
let y = 44;  
let mut p = &x; // L0: x is borrowed
```

```
p = &y; // L1: y is borrowed, overwrites p
```

```
x += 1; // OK, because p was overwritten  
println!(*p); // 44
```

Why?

- Easier development and experimentation
- Common logic *implemented once*
- Similar logic *looks similar*
- Promising Performance in the Literature™
- More advanced analysis



How?

```
let mut x = 22;  
let y = 44;  
let mut p: &'a i32 = &'b x; // L0
```

'b \subseteq 'a
{L0} \subseteq 'b



```
p = &'c y; // L1
```

'c \subseteq 'a
{L1} \subseteq 'c
'a: 0, I am slain! ☠



Prov

```
x += 1;  
println!(*p);
```

invalidates(L0)



How??

```
provvar_live_at(R, P) :-  
    var_drop_live(V, P),  
    var_drops_provvar(V, R).
```

```
provvar_live_at(R, P) :-  
    var_live(V, P),  
    var_uses_provvar(V, R).
```

```
var_live(V, P) :-  
    var_live(V, Q),  
    cfg_edge(P, Q),  
    !var_defined(V, P).
```

```
provvar_live_at.from_join(  
    &var_drop_live,  
    &var_drops_provvar, |_v, &p, &r| {  
        ((r, p), ())  
    });
```

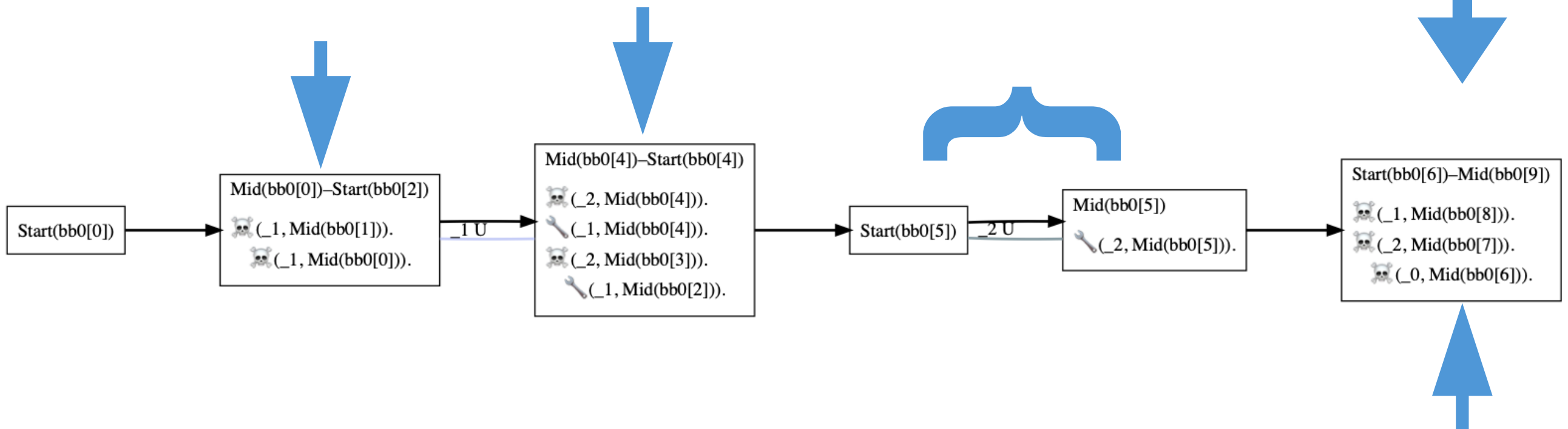
```
provvar_live_at.from_join(  
    &var_live_var,  
    &var_uses_provvar, |_v, &p, &r| {  
        ((r, p), ())  
    });
```

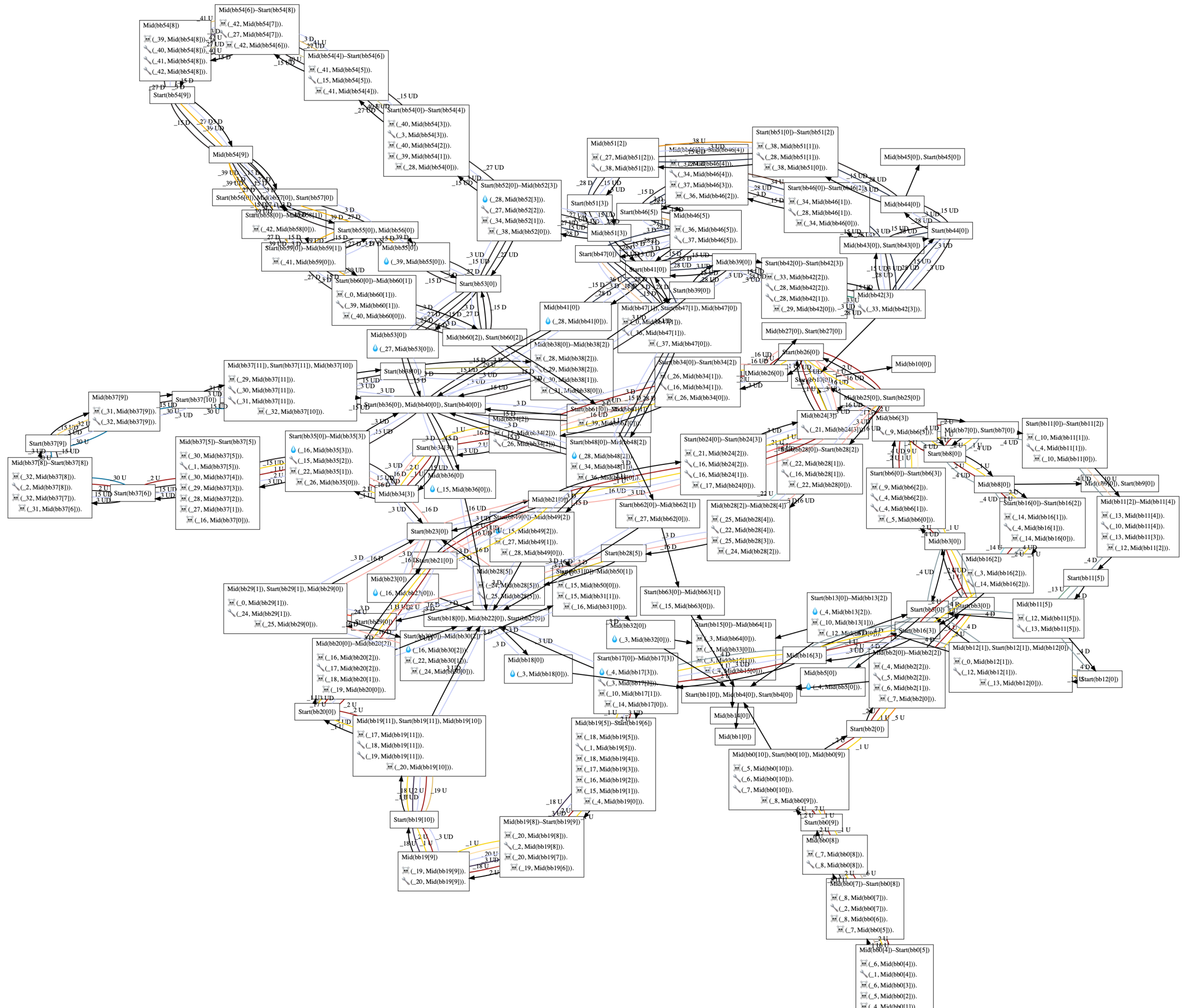
```
var_live_var.from_leapjoin(  
    &var_live_var,  
    (  
        defined_rel.extend_anti(  
            |&(v, _q)| v),  
        edge_reverse_cfg.extend_with(  
            |&(_v, q)| q),  
    ),  
    |&(v, _q), &p| (v, p),  
);
```

How???

```
fn foo() {  
  let x = 22;  
  let p = &x;  
}
```

```
bb0: {  
  { StorageLive(_1);  
    _1 = const 22i32;  
    FakeRead(ForLet, _1);  
  }  
  { StorageLive(_2);  
    _2 = &_1;  
    FakeRead(ForLet, _2);  
  }  
  { _0 = ();  
    StorageDead(_2);  
    StorageDead(_1);  
  }  
  return;  
}
```





What dreams may come?

- (de-)initialisation tracking (also impacts liveness)
- re-formulation of the subset rules upon set equality
- illegal subset relations
- higher-ranked provenances (e.g. `f<'b>_(&'b u32)_`)
- validation
- pre-benchmarking

Summary

- Polonius: a reformulation of the Rust borrow check
- Provenance variables in types = sets of possible source loans
- Declaratively defined in “Datalog”
- Example of prov-var liveness

